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## PATENT ABSTRACTS OF JAPAN

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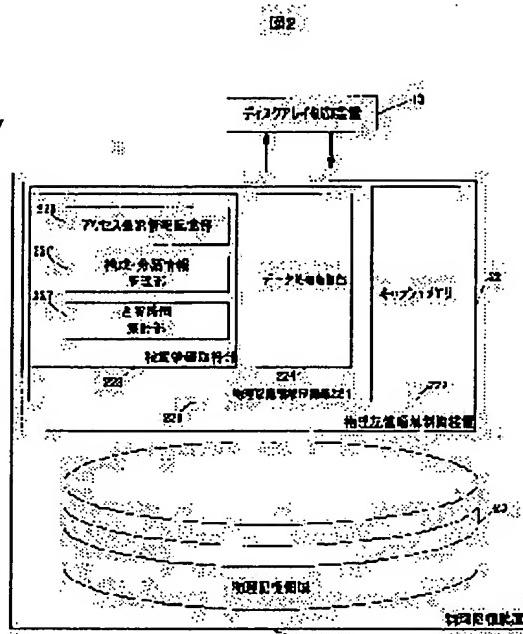
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## (54) STORAGE SUB-SYSTEM, AND MEMORY USED THEREFOR

## (57)Abstract:

**PROBLEM TO BE SOLVED:** To obtain an occupied time of a logic storage area in a physical memory, and to obtain precise access occupied time information in every I/O to the physical memory.

**SOLUTION:** A physical storage area controller 22 on the individual physical memory 15 is provided with a table 225 for storing information about access requirement from a host computer, a table 227 for totalizing the occupied time as to access, a table 226 for control information for classifying constitution of a disk array, and a data processing control part 224 for obtaining constitution information and classification information of the logic storage area form a disk array controller 13, and for requesting the constitution information and the classification information of the logic storage area to the disk array controller, when necessary. The disk array controller 13 is provided with a means for transmitting the constitution information of the disk array at the present time to the physical storage area controller in response to the request from the physical storage area controller on the physical memory.



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## CLAIMS

## [Claim(s)]

[Claim 1] A means to be connected to 1 or two or more computers, and to acquire the operating condition information on two or more physical memory equipments and two or more of these physical memory equipments. In the storage subsystem which has a means to perform matching with the logic storage region which said computer makes a read/write object, and the physical memory field of said physical memory equipment It is the storage subsystem which each of two or more of said physical memory equipments is equipped with a physical memory field control unit, and is characterized by this physical memory field control unit having a means to acquire the operating condition of a physical memory field.

[Claim 2] A means to be connected to 1 or two or more computers, and to acquire the operating condition information on two or more physical memory equipments and two or more of these physical memory equipments. In the storage subsystem which has a means to perform matching with the logic storage region which said computer makes a read/write object, and the physical memory field of said physical memory equipment A means to acquire the operating condition information on said two or more physical memory equipments, and a means to perform matching with the logic storage region which said computer makes a read/write object, and the physical memory field of said physical memory equipment It is prepared in the control unit which controls two or more above-mentioned physical memory equipments. Said control unit Furthermore, it has a means to transmit the information which performed matching with the logic storage region of physical memory equipment, and the physical memory field of physical memory equipment to said two or more physical memory equipments of each. It is the storage subsystem which each of two or more of said physical memory equipments is equipped with a physical memory field control unit, and is characterized by this physical memory field control unit having a means to acquire the operating condition of a physical memory field.

[Claim 3] It is physical memory equipment which is equipped with a physical memory field control unit in the physical memory equipment which constitutes a storage subsystem, and is used for the storage subsystem according to claim 1 or 2 characterized by this physical memory field control unit having a means to acquire the operating condition of a physical memory field.

[Claim 4] Said physical memory field control unit is physical memory equipment according to claim 3 characterized by having further a means to store the operating condition information on the acquired physical memory field.

[Claim 5] Said physical-memory field control unit is physical-memory equipment according to claim 3 or 4 characterized by to have further a means store the information which matched the logic storage region and the physical memory field of the physical-memory equipment which receives the operating condition information on the physical memory field of self-physical memory equipment from a means to transmit to said control unit, and said control unit, according to the acquisition demand of the operating-condition information on a physical-memory field which receives from said control unit.

[Translation done.]

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## DETAILED DESCRIPTION

## [Detailed Description of the Invention]

[0001]

[Field of the Invention] This invention relates to the store used for a storage subsystem and its system, and relates to the store used for the storage subsystem which has two or more stores especially, and its system.

[0002]

[Description of the Prior Art] The disk array system is known as a highly efficient secondary-storage system used for a computer system.

[0003] A disk array system is a system which made it possible to perform read/write of the data which arrange two or more physical memory equipments in the shape of an array, divide and store data in each physical memory equipment, and said each physical memory equipment is operated to juxtaposition, and are stored in said each physical memory equipment by dividing at a high speed.

[0004] As a conventional technique about a disk array system, the technique indicated by D.A.Patterson, G.Gibson, and R.H.Kats, "A Case for Redundant Arrays of Inexpensive Disks (disk array) etc.", (inch Proc.ACM SIGMOD, pp.109-116, and June 1988), etc. is known. This conventional technique gives the classification of level 5 from level 1 to the disk array system which added redundancy according to that configuration. Moreover, a disk array system without redundancy may be added to these classification, and this may be called level 0 to them. In order to realize each above-mentioned level as a different configuration according to redundancy etc., cost differs from performance characteristics etc. And in building a disk array system, the array (group of physical memory equipment) of two or more level is made intermingled in many cases. Here, the group of the disk array which added redundancy is called a parity group. Moreover, in order to realize optimal cost performance in costs changing with the engine performance, capacity, etc. also about physical memory equipment, and building a disk array system, two or more sorts of physical memory equipments with which the engine performance differs from capacity too may be used.

[0005] It distributes to said above physical memory equipments, and the data stored in a disk array system are arranged. For this reason, a disk array system needs to perform matching of the physical memory field which shows the logic storage region which the host computer connected to a disk array system accesses, and the storage region of said physical memory equipment, i.e., address translation.

[0006] The technique indicated by JP9-274544 A etc. is known as a conventional technique about the disk array system which processes address translation. This conventional technique says that a means to acquire the information about I/O access over the logic storage region from a host computer, and a means to change matching with the physical memory field of a logic storage region, and to perform physical relocation realize the optimal arrangement of the stored data, and the technique in which a disk array control unit acquires the I/O access occupancy hour entry to a logic storage region is indicated by this official report.

[0007] In order to perform the load distribution of a disk array system, in case it changes matching with the physical memory field of a logic storage region and performs physical

relocation, since the I/O access occupancy hour entry to a logic storage region serves as data which become origin, it is important.

[0008] [Problem(s) to be Solved by the Invention] Although, as for the conventional technique indicated by the official report mentioned above, it is said that a disk array control unit acquires the occupancy time amount by I/O to a logic storage region, the approach shown in this conventional technique has the trouble that it explains below.

- 2.\*\*\* First, the case where the writing (light) of data is performed to a certain logic storage region is considered. In this case, the light of data is performed to the physical memory field corresponding to the logic storage region concerned. A physical memory field is in physical memory equipment, and physical memory equipment is constituted by the cache memory and the physical memory field which mainly carry out the cache of a physical memory field control section and the data. And when carrying out the light of the data to a physical memory field, it writes in, when a physical memory field control section writes light data in a cache, and the response of termination is notified to a disk array control unit. For this reason, in order that the conventional technique mentioned above may actually carry out the light of the data, it will have the trouble that the time amount which accessed the physical memory field is not known.
3. Next, the case where reading (lead) of data is performed to a certain logic storage region is considered. In this case, in fact, although the lead of the data in the physical memory field corresponding to the logic storage region concerned is performed, when that lead data is in the cache memory in physical memory equipment, a physical memory field is not accessed, but cache memory is accessed, and that data is returned. For this reason, it will have the trouble that the exact time amount which could not distinguish and accessed the physical memory field does not understand whether the conventional technique mentioned above actually had access in the physical memory field for the data lead.

[0011] Moreover, the case where the logic storage region currently crossed to A, B, C, D, and multiple times has access is considered. And the response of 'A' and B is made into 'B' etc. for the response of A, and time of day of B' (t) and response 'A' is made [the time of day of Access A] into A' (t) etc. for the time of day of A (t) and Access B. Here, Access A accesses physical memory equipment with a data lead, leads data, and assumes Access B to be what had the response B without accessing physical memory equipment, since data are in a cache with a data lead. In this case, the direction of response 'B' of Access B which came after response A of the access A which suited previously becomes previously. That is, it becomes A(t) < B(t), A'(t) > B' (t). At this time, the conventional technique which acquires the occupancy time amount by I/O to a logic storage region in a disk array control unit produces the trouble that sweet red bean soup with mochi cannot perform whether which I/O accessed the physical memory field in physical memory equipment how much, or it hit into the cache of physical memory equipment.

[0012] The purpose of this invention is by solving the trouble of the conventional technique mentioned above and acquiring the occupancy time amount of a logic storage region with physical memory equipment to offer the store which offers the storage subsystem which enabled it to acquire the occupancy time amount (real operating time) of each logic storage region by the system configuration unacquirable [with a disk array control unit], and is used for this.

[0013] The purpose of this invention moreover, by taking into consideration the effect on the physical memory equipment for every I/O in order to be able to make small the error of the analysis of the utilization factor of the logic storage region of said disk array system, or utilization factor prediction and to perform optimum performance tuning more it is in offering the store which offers the storage subsystem which can acquire the access occupancy hour entry for every I/O to physical memory equipment with a more high precision, and is used for this.

[Means for Solving the Problem] According to this invention, said purpose is connected to 1 or two or more computers. Two or more physical memory equipments, [in the storage subsystem which has a means to acquire the operating condition information on two or more of these physical memory equipments, and a means to perform matching with the logic storage region which said computer makes a read/write object, and the physical memory field of said physical

memory equipment A means to acquire the operating condition information on said two or more physical memory equipments, and a means to perform matching with the logic storage region which said computer makes a read/write object, and the physical memory field of said physical memory equipment It is prepared in the control unit which controls two or more above-mentioned physical memory equipments. Said control unit Furthermore, it has a means to transmit the information which performed matching with the logic storage region of physical memory equipment, and the physical memory field of physical memory equipment to said two or more physical memory equipments of each. It has a physical memory field control unit, and this physical memory field control unit is attained by having a means of two or more of said physical memory equipments to acquire the operating condition of a physical memory field, respectively. [0015] Moreover, said purpose is set to the physical memory equipment which constitutes a storage subsystem. A means by which have a physical memory field control unit and this physical memory field control unit acquires the operating condition of a physical memory field. An acquisition demand of the operating condition information on a means to store the operating condition information on the acquired physical memory field, and the physical memory field received from said control unit is accepted. It is attained by having a means to transmit the operating condition information on the physical memory field of self-physical memory equipment to said control unit, and a means to store the information which matched the logic storage region and physical memory field of the physical memory equipment received from said control unit.

[Embodiment of the Invention] Hereafter, a drawing explains the operation gestalt of the storage used for the storage subsystem by this invention, and its system to a detail. [0017] The block diagram showing the configuration of the computer system with which drawing 1 was equipped with the storage subsystem by this invention, and drawing 2 are the block diagrams showing the configuration of physical memory equipment. In drawing 1 R> 1 and drawing 2 a host and 12 10 A storage subsystem, 13 disk array control information and 15 for a disk array control device and 14 Physical memory equipment. In 16, a disk array and 17 an I/O bus and 19 for a control terminal and 18 A network, 22 a physical memory field and 130 for a physical memory field control unit and 23 The read/write processing section, 131 the relocation decision processing section and 133 for the operating condition information acquisition processing section and 132 The relocation executive operation section, 141 class configuration information and 143 for the information corresponding to logic/physics, and 142 Class attribute information. 144 physical field operating condition information and 146 for logic field operating condition information and 145 Relocation decision horizon information, 147 free-space information and 149 for relocation activation time information and 148 Relocation information, 14A --- a storage occupancy hour entry and 221 --- for the operation information acquisition section and 224, as for the access request information storage section and 226, a data-processing control section and 225 are [ a physical memory field control section and 226 / cache memory and 223 / configuration / classification Research and Data Processing Department and 227 ] the occupancy time amount total sections.

[0018] The computer system shown in drawing 1 consists of 1 which is the calculating machine of a high order or two or more hosts 10, a storage subsystem 12, and a control terminal 17. It connects with the storage subsystem 12 by I/O bus 18, and a host 10 publishes lead of data, and I/O for light processing to the storage subsystem 12. In case this I/O is performed, a host 10 specifies the logical storage region of the storage subsystem 12. That is, a host 10 accesses with the address of a usually logical storage region to the data in a storage subsystem. Moreover, I/O bus 18 is constituted by ESCON, SCSI, the fiber channel etc.

[0019] The storage subsystem 12 consists of a disk array control unit 13 is equipped with the read/write physical memory equipments 15. The disk array control unit 13 is equipped with the read/write processing section 130, the operating condition information acquisition processing section 131, the relocation decision processing section 132, and the relocation executive operation section 133, and these processing sections process read/write processing, operating condition information acquisition processing, relocation decision processing, relocation executive operation, etc. Moreover, the disk array control non-dense of the storage subsystem 12 holds the

information 141 corresponding to a logic storage region / physical memory field, the class configuration information 142, the disk array configuration information 1400 of class attribute information 143 grade, storage occupancy hour entry 14A of the logic field operating condition information 144 and physical field operating condition information 145 grade, the relocation decision horizon information 146, the relocation activation time information 147, the free-space information 148, and relocation information 149 grade. In addition, others and parity group information, RAID level information, etc. which were mentioned above may be included in the disk array configuration information 14 mentioned above. [ information ] [0020] Moreover, the host 10, the disk array control unit 13, and the control terminal 17 are mutually connected by the network 19. A network 19 may be constituted by Ethernet (trademark), FDDI, the fiber channel, etc. A control terminal 17 is usually used in order to perform maintenance, management, etc. of the storage subsystem 12. [0021] Moreover, in explanation of the operation gestalt of this invention, although it exists, respectively, since it is not important, the component which surely exists in computers, such as memory for performing processing which appears in a host 10, the disk array control unit 13, and a control terminal 17, respectively, and CPU, is not specified here.

[0022] The class division of two or more physical memory equipments 15 formed in the above-mentioned storage subsystem 12 is carried out for every engine performance of physical memory equipment, and they constitute the disk array 16 for every class. Moreover, although not shown clearly, two or more physical memory equipments are used, and the parity group is constituted here. And each of physical memory equipment 15 is constituted by the physical memory field control device 22, which controls the physical memory field 23 and this physical memory field 23 to be shown in drawing 2, and various data are stored in the physical memory field 23. [0023] Moreover, if the address of the physical memory field 23 seems to have mentioned above directly from a host 10, it does not break, but a host 10 accesses the data on two or more logical storage regions on two or more physical memory fields 23. Namely, a host 10 accesses by specifying a logic storage region to the data in the storage region of each physical memory equipment 15 in the storage subsystem 12. [0024] The disk array control unit 13 is connected with two or more physical memory equipments 15. The lead and light processing instruction I/O which controlled two or more physical memory equipments 15, or were emitted by said host 10 Make the address of the logic storage region where the appointed data exist, and the address of a physical memory field with the address of the logic storage region correspond, transmit data I/O to suitable physical memory equipment 15, and if it is light processing, The data transmitted by the host 10 are transmitted to physical memory equipment 15, and if it is lead processing, the data transmitted from physical memory equipment 15 are received, and it is processing transmitting to a host 10 etc.

[0025] The physical memory field control unit 22 which it has in physical memory equipment 15 is constituted by the physical memory field control section 221 and cache memory 222. Cache memory 222 has the quick rate of processing of the read/write of data compared with the physical memory field 23. And cache memory 222 is used as follows about the data about the lead or light instruction transmitted from the disk array control device 13. That is, in case the light data transmitted from the disk array control device 13 are written in the physical memory field 23 in light processing, data are written also in cache memory 222. Moreover, in case data reading appearance is carried out from the physical memory field 23 in lead processing, when it was written in cache memory 222, or the same data are in cache memory and it comes to the physical memory equipment from the disk array control device 13 as a lead instruction to the data by former lead processing, the read data do not read the data from the physical memory field 23, but read it from cache memory 222. Thereby, the processing engine performance of physical memory equipment 15 can be improved.

[0026] The physical memory field control section 221 is mainly equipped with the operation information acquisition section 223 and the data-processing control section 224, and is constituted. The data-processing control section 224 receives the lead or light instruction of data transmitted from the disk array control device 13. And the data-processing control section

224 will read the lead data from cache memory 222, if cache memory 222 is accessed and the lead data exists in cache memory 222, when the received instruction is a lead instruction, if the data is not in cache memory 222, it will access the physical memory field 23, will read the lead data, and will transmit data to the disk array control device 13. Moreover, the data-processing control section 224 is after that, and writes the data in the physical memory field 23 at the same time it writes the data transmitted from the disk array control device 13 in cache memory 222, when the received instruction is a light processing instruction. You may not write light data in cache memory 222, but may also write them in the direct physical memory field 23.

[0027] The operation information acquisition section 223 is constituted by the access request information storage section 225, configuration / classification Research and Data Processing Department 226, and occupancy time amount total section 227 grade. When a logic storage region with the data specified by the I/O process, the physical memory field 23 where the data exists, or cache memory 222 is accessed, the above-mentioned data-processing control section 224 classifies the hour entry of the access for every (random access, sequential access, etc.) processing classification of an I/O process, and records it on the occupancy time amount total section 227. Moreover, from the disk array control device 13, the data-processing control section 224 receives information, such as correspondence information on the address of the logic storage region in a disk array, and the address of the physical memory field 23, and engine performance of physical memory equipment 15, and records it on configuration / classification Research and Data Processing Department 226. Furthermore, the data-processing control section 224 receives a lead of data or light instruction data transmitted from disk array control device 13 grade, and records the instruction data on the access request information storage section 225 so that it may be possible to receive two or more I/O processes.

[0028] The logic storage region according to an I/O process within physical memory equipment 15 when physical memory equipment 15 has a configuration which was mentioned above, the physical memory field 23. Or it becomes possible to classify the hour entry at the time of access when accessing cache memory 222 for every (random access, sequential access, etc.) processing classification of an I/O process, and to record on the occupancy time amount total section 227. It classifies into a physical memory field to what whether time amount access was carried out and physical memory equipment 15 of cache memory 22 it hit according to an I/O process, and it becomes possible to total the occupancy time amount.

[0029] Drawing 3 is a flow chart explaining processing actuation of a disk array control device when a storage subsystem is started, and explains this hereafter.

[0030] (1) The disk array control unit 13 transmits the information 141 corresponding to the logic/physics which is the matching information on the address of the logic storage region in the physical memory field 23 in physical memory equipment 15, and the address of the physical memory field where the logic storage region actually exists, the class configuration information 142, and the disk array configuration information 14 of class attribute information 143 grade to the physical memory equipment 15 connected with self-equipment 13 at the time of starting of the storage subsystem 12 (steps 300 and 310).

[0031] (2) Next, the disk array control unit 13 receives the notice the physical memory equipment 15 sent from the physical memory field control unit 22 changed [the notice] to the accessible ready state, when physical memory equipment 15 becomes accessible by transmission of the information mentioned above. Things come and physical memory equipment 15 is in the condition of the initialization termination by the disk array configuration information 14 (step 320).

[0032] (3) Then, the disk array control device 13 receives data with more various what has transmitted host I/O of a lead or light processing to the storage subsystem 12 to the logic storage region in the storage subsystem 12, things which deliver an instruction and data with disk array control devices than a host 10 by I/O-bus 18 course (step 330).

[0033] (4) When host I/O is received as said received data, the disk array control unit 13 receives the lead or light demand to the logic storage region specified by host I/O, and asks for the logic storage region address and the corresponding address of the physical memory field 23 using the information 141 corresponding to the logic/physics which changes the address (logical

address) of the logic storage region into the address (physical address) of a physical memory field (steps 340 and 350).

[0034] (5) The disk array control device 13 specifies the address of the physical-memory field where predetermined data exist, and, in lead processing, transmits light data to the physical-memory equipment which transmits lead data for the physical memory equipment which has the above-mentioned physical address to lead data to read-out and a host 10, receives the light data transmitted by the host 10 in light processing, and has the physical address (360).

[0035] Drawing 4 is a flow chart explaining processing actuation of a disk array control device when correspondence of the address of a logic storage region and the address of a physical memory field changes, and explains this hereafter.

[0036] (1) By the change of RAID level and a logic storage region moving the disk array control device 13 to the address of a physical memory field different from a certain physical memory \*\*\* address now etc. When it supervises that correspondence of the address of a logic storage region and the address of a physical memory field changed and is changeful, physical memory equipment 15 is received again. The information 141 corresponding to the logic/physics which is the matching information on the address of the logic storage region in the physical memory field 23 in physical memory equipment 15, and the address of the physical memory field where the logic storage region actually exists. The class configuration information 142 and the disk array configuration information 14 of class attribute information 143 grade are transmitted (step 310).

[0037] (2) Next, the disk array control unit 13 receives the notice the physical memory equipment 15 sent from the physical memory field control unit 22 changed [the notice] to the accessible ready state, when physical memory equipment 15 becomes accessible by transmission of the information mentioned above. Things come and physical memory equipment 15 is in the condition of the updating termination by the disk array configuration information 14 (step 320).

[0038] (3) Subsequent processing is performed like the case of steps 330, 340, 3500, and 3600 explained by drawing 3 (steps 3301, 3401, 3501, and 3601).

[0039] In addition, in steps 320 and 3201 mentioned above, as for the disk array control unit 13, \*\* does not need to receive the information that the physical memory field control unit 22 to physical memory equipment 15 changed in the accessible condition. In this case, what is necessary is just to perform access processing for performing predetermined processing to physical memory equipment 15, when it assumes that it is in accessible conditions, such as a lead and a light. To physical memory equipment 15 and the access directions to a certain storage region come for the disk array control unit 13 after a certain fixed time amount. Moreover, when there is no response into waiting and fixed time amount until it performs access processing in order to perform predetermined processing again when there is no response to predetermined access processing, or a response comes back, it is good also as a method which tells that to the module which issued the access directions to a certain storage region.

[0040] Moreover, the information 141 corresponding to logic/physics in the above-mentioned is information to which a logic storage region and a physical memory field are made to correspond. And the logical address is the address which shows the logic storage region which a host 10 uses in said read/write processing section 130. Moreover, a physical address is the address which shows the field on the physical memory equipment 15 with which data are actually stored, and consists of the physical memory device number and the address in physical memory equipment. A storage number shows each physical memory equipment 15. The address in storage is the address which shows the storage region within physical memory equipment 15.

[0041] Drawing 5 is a flow chart explaining processing actuation of the disk array control device 13 at the time of the disk array control device 13 reading the information in the operation information acquisition section 223 of physical memory equipment 15, and explains this hereafter.

[0042] (1) After the storage subsystem 12 is started, the disk array control device 13 initializes storage occupancy hour entry 14A, and transmits an acquisition demand of the access occupancy hour entry of the physical memory equipment 15 to two or more physical memory

equipments 15 connected after that (steps 371 and 372).

[0043] (2) Next, the disk array control unit 13 stores the access occupancy hour entry of reception and each physical memory equipment in storage occupancy hour entry 14A for an access occupancy hour entry from each physical memory equipment 15 (steps 373 and 374).

[0044] In addition, the timing of acquisition of the access occupancy information on the disk array control unit 13 mentioned above The method which reads the access occupancy information by access to physical memory equipment 23 from the module of others by host I/O, backup, etc. to a fixed time interval from the occupancy time amount total section 227 in each physical memory equipment 15. When an access occupancy hour entry acquisition demand is transmitted to the disk array control unit 13 from other modules (for example, a host 10 and a control terminal 17), it is various and dependent on a design.

[0045] The access occupancy hour entry acquired by the above-mentioned is recorded on the occupancy time amount total table in the disk array control device 13.

[0046] Drawing 6 is a flow chart explaining processing actuation of the physical memory field control device 22 in physical memory equipment 15, next explains this.

[0047] (1) The physical memory field control device 22 receives the disk array configuration information 14 which is information on the logical address and the physical address in the physical memory field 23 of the data 15 transmitted from the disk array control device 13 at the time of starting of the storage subsystem 12, i.e., physical memory equipment, such as correspondence, (steps 400 and 401).

[0048] (2) The physical memory field control unit 22 which received the disk array configuration information 14 performs creation initialization of configuration / classification information management table of configuration / classification Research and Data Processing Department 226 in the operation information acquisition section 223, or the occupancy time amount total table of the occupancy time amount total section 227 based on the information (step 402).

[0049] (3) After initialization processing of the table in step 402 is completed, in order to make this physical memory equipment 15 recognize an accessible thing, notify that initialization processing was completed to the disk array control unit 13. In addition, it is not necessary to transmit the information to which physical disk equipment changed in the accessible condition to the disk array control device 13. When the access directions to a certain storage region come to the physical memory field control unit 22 of physical memory equipment 15 after the fixed time amount, in this case, the physical memory field control unit 22 If it is in a condition accessible to the physical memory field 23 in order to perform predetermined processing Access the physical memory field 23, perform predetermined processing, and if impossible [ whether processing predetermined in the condition of having stored access information in the access request information storage section 225, and having become accessible to the physical memory field is performed, and ] or you may make it not receive a certain access directions to a storage region until it comes to resemble the physical memory field 23 in the accessible condition, in order to perform predetermined processing (step 403)

[0050] (4) After that, the physical memory field control device 22 waits to transmit host I/O, and a physical memory equipment operation information acquisition demand instruction or new disk array configuration information from the disk array control device 13, and receives it (step 404).

[0051] (5) At step 404, if host I/O is received from the disk array control unit 13, when the I/O judges lead processing or light processing and it is lead processing, the data processing control section 224 will confirm whether the data which should be read exist in cache memory 22 (steps 405 and 406).

[0052] (6) When the data is read from cache memory 22 when the data exists in cache memory 22, and the data does not exist in cache memory 22 with the check of step 406, read the data from the physical memory field 23, and transmit data to the disk array control unit 13 (steps 407, 709, and 408).

[0053] (7) At step 405, when judged with host I/O being light processing, the data-processing control section 224 receives the light data transmitted by the host 10, and writes the light data in cache memory 22 (steps 410 and 411).

[0054] (8) And the data-processing control section 224 stores the above-mentioned light data in time amount total table in the receiving disk array control device 13. The example shown in

the physical memory field 23 while notifying the notice of data write-in termination to the disk array control device 13 (steps 412 and 413).

[0055] The data-processing control section 224 to cache memory 22 after processing of step 408, or processing of step 413 (9) In access \*\* Or the JOB classification information on the information on whether the physical memory field 23 was accessed, a random lead, a sequential lead, etc.. The access classification of the JOB classification information on the random read/write at the time of writing light data in the physical memory field 23, sequential read/write, etc. is recognized. The occupancy hour entry which accessed cache memory 22 or the physical memory field 23 is stored in the occupancy time amount total section 227 in the operation information acquisition section 223 for every access classification (steps 414 and 415).

[0056] (10) If new disk array configuration information is received from the disk array control

device 13 at step 404, the data-processing control section 224 will rewrite the information in configuration / classification Research and Data Processing Department 226, the operation

information acquisition section 223, corresponding to new disk array configuration information (steps 418 and 419).

[0057] (11) When a physical memory equipment operation information acquisition demand instruction is received from the disk array control device 13 at step 404, the data-processing control section 224 reads the access occupancy hour entry of the physical memory equipment 15 stored in the occupancy time amount total section 227 in the operation information acquisition section 223, and transmits it to the disk array control device 13 (steps 416 and 417). [0058] In addition, it may be made to perform transmission to the disk array control device 13 from the physical memory equipment 15 of the occupancy hour entry in processing of step 417 mentioned above to the disk array control device 13 with a fixed time interval automatically from physical memory equipment 15. In this case, said physical memory equipment operation information acquisition demand instruction is not transmitted to physical memory equipment 13 from the disk array control unit 13.

[0059] Drawing 7 is drawing explaining the example of a configuration of the information 141 corresponding to logic/physics in the table for managing correspondence with the address of a logic storage region and the address of a physical memory field which are held in the disk array control unit 13.

[0060] The disk array control unit 13 has managed correspondence with the address of the logic storage region in the physical memory field 23 in two or more physical memory equipments 15 connected, and the address of the physical memory field in the logic storage region. Each of the logic storage region number 500 given to a specific logic storage region, the logical address 510, the physical address 520 by the storage number 521 with a physical storage region with the logic storage region and the address 522 of a physical storage region, the RAID level 530 that shows the engine performance of the physical-memory equipment 15, and the parity group number 540 to which the physical-memory equipment 15 belongs is matched, and the information 141 corresponding to the logic/physics used for this is constituted, as shown in drawing 7. When processing of a lead, a light, etc. specifies the address of a logic storage region from a host 10, a control terminal 17, and other modules (for example, other disk array control units etc.) and it is accessed to the self-disk array control unit 13 by having such information 141 corresponding to logic/physics, the disk array control unit 13 can change the address of a logic field into the address of a physical storage region, and can perform read/write processing of data correctly to physical memory equipment 15.

[0061] Drawing 8 is drawing showing the example of the logic field operating condition information 144 stored in the disk array control unit 13, and the storage occupancy hour entry 141 of physical field operating condition information 145 grade. Such information is constituted as an occupancy time amount total table.

[0062] The disk array control device 13 reads periodically the access occupancy information on the physical memory field 23 of the physical memory equipment 15 by access from the module of others by host I/O, backup, etc. from the occupancy time amount total section 227 in each physical memory equipment 15, and records the access occupancy hour entry on the occupancy time amount total table in the receiving disk array control device 13. The example shown in

drawing 8 is every logic storage region number 601 and I/O. Occupancy time amount is totaled every JOB classification 602. I/O As a JOB classification, although 670 and a total of 880 are shown by the example of illustration at the time of the sequential lead 610, the sequential light data 620, the sequential light parity 630, the random lead 640, the random light parity 680, and a cache hit, it is I/O of further others. There may be JOB classification.

[0063] The occupancy time amount read from physical memory equipment 15 may be record by the accumulation value of not only the above-mentioned but the access occupancy time amount for every I/O, and a universal time amount value and a time amount value peculiar to a machine. Moreover, in the disk array control unit 13, the access occupancy time amount of each logic storage region or a physical memory field may be edited, and an occupancy hour entry table may newly be created based on the value which found the access occupancy time amount for every physical memory equipment and every parity group.

[0064] When a host 10 and control terminal 17 grade give the operation information acquisition demand of physical memory equipment to the disk array subsystem 12 by the above-mentioned even if a host 10 and control terminal 17 grade access direct physical memory equipment 15 and do not acquire the access occupancy hour entry of a logic storage region or a physical memory field, they become possible [ acquiring the access occupancy hour entry of the logic storage region and physical memory field from the disk array control unit 13 ].

[0065] Drawing 9 is drawing showing the example of a configuration of the table which manages

matching with the address of a logic storage region and the address of a physical memory field

which are stored in configuration / classification Research and Data Processing Department 226

in the operation information acquisition section 223 in physical memory equipment 15.

[0066] From the disk array control device 13, when the relation of the address of a logic storage region and the address of a physical-memory field which are produced by the time of starting of the storage subsystem 12, the change in physical-memory equipment, change of RAID level, migration of a logic storage region, etc. changes, the data-processing control section 224 in physical-memory equipment 15 receives the matching information on the logic storage region and physical-memory field, and stores it in a correspondence table as shows it to configuration / classification Research and Data Processing Department 226 at drawing 9. This correspondence table is constituted by physical ADDRE 7520 by the logic storage region number 700 given to a specific logic storage region, the logical address 710, and the storage number 721 with a physical storage region with that logic storage region and the address 722 of a physical storage region. Thereby, it can recognize which address logic storage region physical memory equipment 15 has where of physical memory regional address in self-physical memory equipment 15.

[0067] Drawing 10 is drawing showing the example of a configuration of the table of the occupancy hour entry by access to the storage region by I/O accumulated and stored at the occupancy time amount total section 227 of the operation information acquisition section 223 in physical memory equipment 15.

[0068] This table is every logic storage region number 801 and I/O. Occupancy time amount is totaled every JOB classification 802. I/O As a JOB classification, in the example of illustration, although 870 and a total of 880 are shown at the time of the sequential lead 810, the sequential light data 820, the sequential light parity 830, the random lead 840, the random light parity 860, and a cache hit, there may be I/OJOB classification of further others.

[0069] If the data-processing control section 224 in physical memory equipment 15 has access in a store by I/O from a host etc., it will accumulate the occupancy time amount 890 in the occupancy time amount total section 227 every JOB classification 802 of access about each logic storage region 801 with access. This becomes possible to obtain the relation between the number of the logic storage region in physical memory equipment 15, the occupancy time amount by the access classification to the logic storage region, and the sum total occupancy time amount within a certain time amount within physical memory equipment 15.

[0070] According to the operation gestalt of this invention mentioned above, the occupancy hour entry of the storage for every I/O can be acquired, and acquisition of the occupancy hour entry of the storage by access to a storage region can be realized within physical memory equipment.

[0071] According to the operation gestalt of this invention, the thing of two or more physical

memory equipments which constitute a storage subsystem for which the occupancy time amount of the logic storage region of physical memory equipment is acquired becomes respectively possible by the above-mentioned, and the thing of each logic storage region to do for occupancy time amount (real operating time) acquisition becomes possible by the system configuration unacquirable [ with a disk array control unit ]. Moreover, in order according to the operation gestalt of this invention mentioned above to be able to make small the error of the analysis of the utilization factor of the logic storage region of said storage subsystem, or utilization factor prediction and to perform optimum performance tuning more by taking into consideration the effect on the physical memory equipment for every I/O, the access occupancy hour entry for every I/O to physical memory equipment with a more high precision is acquirable.

[0072] [Effect of the Invention] As explained above, according to this invention, physical memory equipment can acquire the occupancy time amount of a logic storage region, and the occupancy time amount (real operating time) of each logic storage region can be acquired by the system configuration unacquirable [ with a disk array control unit ].

[0073] Moreover, according to this invention, since the access occupancy hour entry for every I/O to physical memory equipment is acquirable, it becomes possible to perform the analysis of the utilization factor of the logic storage region of a storage subsystem and the prediction of a utilization factor in consideration of the effect on the physical memory equipment for every I/O with a small error, and the engine performance of the more nearly optimal storage subsystem can be tuned up.

[Translation done]

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1.This document has been translated by computer. So the translation may not reflect the original precisely.

2.\*\*\*\* shows the word which can not be translated.

3.In the drawings, any words are not translated.

## DESCRIPTION OF DRAWINGS

## [Brief Description of the Drawings]

[Drawing 1] It is the block diagram showing the configuration of the computer system equipped with the storage subsystem by this invention.

[Drawing 2] It is the block diagram showing the configuration of physical memory equipment.

[Drawing 3] It is a flow chart explaining processing actuation of a disk array control device when a storage subsystem is started.

[Drawing 4] It is a flow chart explaining processing actuation of a disk array control device when correspondence of the address of a logic storage region and the address of a physical memory field changes.

[Drawing 5] It is a flow chart explaining processing actuation of the disk array control device at the time of a disk array control device reading the information on operation information acquisition circles of physical memory equipment.

[Drawing 6] It is a flow chart explaining processing actuation of the physical memory field control device in physical memory equipment.

[Drawing 7] It is drawing explaining the example of a configuration of the information corresponding to the logic/physics which manages correspondence of the address of the logic storage region and physical memory field which are held in the disk array control unit.

[Drawing 8] It is drawing showing the example of storage occupancy hour entries, such as logic field operating condition information stored in a disk array control unit, and physical field operating condition information.

[Drawing 9] It is drawing showing the example of a configuration of the table which manages matching of the address of the logic storage region and physical memory field which are stored in configuration / classification Research and Data Processing Department in physical memory equipment.

[Drawing 10] It is drawing showing the example of a configuration of the table of the occupancy hour entry to the storage region accumulated and stored at the occupancy time amount total section of the operation information acquisition section in physical memory equipment.

## [Description of Notations]

10 Host

12 Storage Subsystem

13 Disk Array Control Unit

14 Disk Array Control Information

15 Physical Memory Equipment

16 Disk Array

17 Control Terminal

18 I/O Bus

19 Network

22 Physical Memory Field Control Unit

23 Physical Memory Field

130 Read/write Processing Section

## 131 Operating Condition Information Acquisition Processing Section

132 Relocation Decision Processing Section

133 Relocation Executive Operation Section

141 Information corresponding to Logic/Physics

142 Class Configuration Information

143 Class Attribute Information

144 Logic Field Operating Condition Information

145 Physical Field Operating Condition Information

146 Relocation Decision Horizon Information

147 Relocation Activation Time Information

148 Free-Space Information

149 Relocation Information

14A Storage occupancy hour entry

221 Physical Memory Field Control Section

222 Cache Memory

223 Operation Information Acquisition Section

224 Data-Processing Control Section

225 Access Request Information Storage Section

226 Configuration / Classification Research and Data Processing Department

227 Occupancy Time Amount Total Section

[Translation done.]

119) 日本国税課題 (1P) (12) 公開特許公報 (A)

「新物种」

〔從來の技術〕コンピューターシステムに使用される高性  
能二次記憶システムとして、ディスクアレイシステム

〔0003〕ディスクレイシステムは、複数の物理記録技術をアレイ状に配置し、各物理記録技術間にデータを並列的に分割して格納しておき、前記各物理記録技術間に分割して格納されるデータを並列的に前記各物理記録技術間に分割して格納するデータへ前記各物理記録技術間に分割して格納されるデータを高速に移すことを可能としたシステムである。

た。この從事技術は、元々性能を付加したディスクアレイシステムに對して、その構成に応じてレベル1からレベル4までの階層別に、冗長性無しのディスクアレイシステムである。前述の各レベル

た。この從事技術は、冗長性等に応じて異なる構成として実現するため、コストや性能特性等が異なる。そして、ディスクアレイの構成と、それをレベル1からレベル4までの階層別に、冗長性無しのディスクアレイシステムを構築するにあたって世数のレベルのアレイ構成の問題を、元々性能を付加したディスクアレイの用を活用せることも多い。ここで物理記憶装置の用)を活用せることも多い。ここで

データベース等によりコストが異なり、ディスクアレイシステムを採用するにあたってはコストパフォーマンスを最優先する目的で、やはり性能や容量の質なる性質の物理記録装置を用いることがある。

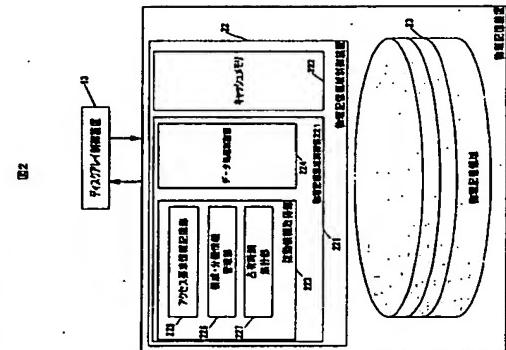
アレイシステムに接続されるホストコンピュータがアクティベーションエコスする論理記憶領域と前記物理記憶装置の記憶領域との対応付け、すなわち、アドレス交換を行う必要がある。

【0006】アドレス交換の処理を行うスクアレイシステムに関する発明技術として、例えば、特開平9-273274号公報等で記載された技術が知られており、この発明技術は、ホストコンピュータからの論理記憶装置に対する1/0アクセスについての情報を取得する手段と、論理記憶領域への対応付けを行う手段により、格別された变形によって物理的再配置を行う手段により、格別された

データの監視機能を実現するというものであり、この公報には、論理記憶領域へのノーアクセス占有時間情報と、ディスクアレイ制御装置が取得する技術が開示されている。

〔0007〕論理記憶領域への1/0アクセス占有時間

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C 0 6 F		3/06	3 0 2 A 3 0 1 X 3 0 4 N 5 4 0 3 2 0 L	7-73-D(参考) 3 0 2 A 3 0 1 X 3 0 4 N 5 4 0 3 2 0 L
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(22)出願日				



## 【発明の名前】ストレージシステム及びそのシステムに使用する記録装置

### 5.7) [要約] **〔採用〕 物理記述装置において、論理記述領域の占有**

ノイズ抑制装置により、物理記憶装置上にある物理記憶装置の初期接続部からリクエストに応じて、現在のディスクアレイの構成情報を物理記憶装置に送信する手段を設ける。



読み出されたデータは、キャッシュメモリ 2/2に書き込まれ、あるいは、以前のリード処理によって同一のデータがキャッシュメモリにあり、そのデータにに対してディスクアレイ制御装置 1 からリード命令として物理記憶装置 2/3から読み込まれ、そこでデータを物理記憶装置 2/3から読み出さず、キャッシュメモリ 2/2から読み込む。これにより、物理記憶装置 1 の処理性能を上げることができる。

【0041】図 5 はディスクアレイ制御装置図 1-3 が物理記憶装置 1-5 の移動情報取扱部 2-2-3 内の情報を出す際のディスクアレイ制御装置図 1-3 の処理動作を説明する。以下、これについて説明する。

【0036】(1) ディスクアレイ制御装置図 1-3 は、物理記憶装置 1-5 の増減や RAID レベルの変化、論理記憶領域が現在ある物理記憶領域アドレスとは別の物理記憶領域のアドレスに移動する等によって、論理記憶領域のアドレスと物理記憶装置 1-5 との接続が変化したことを監視し、変化があった場合、再度、物理記憶装置 1-5 に針對して、物理記憶装置 1-5 内の物理記憶領域アドレスと実際にその論理記憶領域にある論理記憶領域アドレスとの対応が何れかで異なる物理記憶領域アドレスと異なる物理記憶領域アドレスとの対応が何れかである論理/物理記憶領域アドレスと構成情報 1-4 がある場合に、物理記憶装置 1-4 からクラス構成情報 1-4 を送信する(ステップ 3-10)。

【0037】(2) 次に、ディスクアレイ制御装置図 1-3 は、前述した情報の送信により、物理記憶装置 1-5 がアクセス不可能となつたときに、物理記憶装置 1-5 がアクセス可能なレディ状態に遷移した通知を受信する。このとき、物理記憶装置 1-5 は、ディスクアレイ構成情報 1-4 による更新終了の状況となつていい(ステップ 3-20)。

【0038】(3) その後の処理は、図 3 により説明したステップ 3-3-0-0、3-3-0-1、3-3-0-2、3-3-0-3、3-3-0-4 のように順次行なわざる「カウント」<sup>1</sup>が 2-2-1 に送信された場合がアタリで、アタリを検出する。

1. 3501、3601)。	【0045】前述により駆動されたアクセス占有時間情報は、ディスクアレイ制御装置 1 3 内の占時専用データーブルに記録される。
【0039】なお、前述したステップ 3 20、3 220 において、ディスクアレイ制御装置 1 3 は、が物理記憶領域制御装置 2 から、物理記憶装置 1 がアクセス可能状況に遷移したという情報を受信しなくてよい。こ	【0046】図 6 は物理記憶装置 1 5 内の物理記憶領域制御装置 2 2 の処理動作を説明するフローチャートである。これについて説明する。
の場合は、またまとまった時間内に物理記憶装置 1 5 に対してリードやライト等のアクセス可能な状態に対する、何らかの記憶領域へのアクセス指示がディスクアレイ制御装置 1 3 にきた場合に、物理記憶装置 1 5 に所定の処理を行うためのアクセス処理を行えばよい。	【0047】(1) 物理記憶領域制御装置 2 2 は、ストレージシステム 1 2 の初期時にディスクアレイ制御装置 1 3 および記憶装置 2 3 にある論理アドレスと、すなわち、物理記憶装置 1 5 の方が必要なデータ、ドレスとの対応等の情報であるディスクアレイ構成情報 1 4 を受信する(ステップ 4 0、4 01)。
また、所定のアクセス処理に対する応答がない場合、再度、所定の処理を行うためアクセス処理を行うか、あるいは、応答が帰ってくるまで待ち、一定時間中に応答がない場合、何らかの記憶領域へのアクセス指示を出したモジュールに対してその旨を伝える方式としてもよい。	【0048】(2) ディスクアレイ構成情報 1 4 を受信した物理記憶領域制御装置 2 2 は、その情報を元に、構成情報取扱部 2 2 3 内の構成・分担情報管理部 2 2 6 の構成・分担情報管理テーブルや占有時間算出部 2 2 7 の占有時間算出テーブルの作成初期化を行う(ステップ 4 1)。
【0040】また、前述における論理/物理対応情報 1 4 は、論理記憶領域と物理記憶領域とを対応させる情報である。そして、論理アドレスは、ホスト 1 0 が前記リード/ライト処理部 1 3 で用いる論理記憶領域とデータアドレスである。また、物理アドレスは、実際にデータが格納される物理記憶装置 1 5 上の領域を示すアドレスであり、物理記憶装置番号及び物理記憶領域内アドレスからなる。記憶装置番号は、個々の物理記憶装置 1 5 を示す。記憶装置内アドレスは、物理記憶装置 1 5 内での記憶領域を示すアドレスである。	【0049】(3) ステップ 4 0 2 でのテーブルの初期化処理が終ると、この物理記憶装置 1 5 にアクセス可能であることを認識したためディスクアレイ制御装置 1 3 に初期化処理が終ったことを通知する。なお、ディスクアレイ制御装置 1 3 に物理記憶装置 1 5 がアクセス可能な状態に遷移した情報を送信しないてもよい。この場合、定まった時間後に物理記憶装置 1 5 の物理記

【0056】(10)ステップ4.04でディスクアレイ制御装置1.3から新たなディスクアレイ構成情報を受信すると、データ処理部回路2.2.4は、復旧情報部回路2.2.3の構成・分類情報を部回路2.2.6内の情報を削除して新たにディスクアレイ構成情報を対応して書き換える(ステップ4.1.9)。

【0062】ディスクアレイ制御装置13は、ホスト1/0やバックアップ等によるその他のモジュールからのアクセスによる物理記憶装置15の物理記憶領域23へのアクセス占有情報を各物理記憶装置15内の占有時間集計部227から定期的に読み出して、そのアクセス占有時間情報を受信ディスクアレイ制御装置11内の占有

(0050) (4) その後、物理記化銀刷印面接写真 2  
 ノゾムヘ引にによって次第に(0057)の整理を行つか、あるいは  
 は、所定の処理を行うために物理記化銀刷印面接写真を用いて  
 は、可視能性度になるにあたる。アセ  
 セス指示を受け付けるようにしてよい(ステップ 4  
 0 3)。

(0057) (1) (II)ステップ 4 0 4 でディスクアレイ制  
 銀接写真 1 3 から物理記化銀刷印面接写真を要求命令を受  
 けた場合、データ処理別用回路 2 3 にアセ  
 ル部 2 3 内の占有時間規制回路 2 7 に格納している物理  
 記憶装置 1 5 のアクセス占有時間情報を読み出して、そ  
 れをディスクアレイ制銀刷印面接写真 1 3 に送信する(ステップ

時間集計テーブルに記録する。図8に示す例は、論理記憶領域番号601毎、1／0 JOB種別602毎に占有時間を集計したものである。1／0 JOB種別として、図示例では、シーケンシャルリード610、シークエンシャルライト6120、シークエンシャルライトバティ630、ランダムリード640、ランダムライトバティ6430、ランダムライト6440の6種類である。

【00631】物理配達装置 5 から読み出される占有時間は、前述に限らず、1/O毎のアカセス占有時間の累積値や、ユニバーサルな時間間隔、マシン固有な時間間隔ト等のアカセス時間間隔を示す。

【0059】図7はディスクアレイ制御装置1・3内に保持されている論理配信領域のアドレスと物理記憶領域のアドレスとの対応を管理するためのテーブル内の論理アドレスとの構成を説明する図である。

【0060】ディスクアレイ制御装置1・3は、接続され

20 装置1・3に送信されることはない。

21 あつた場合、データ処理部2・2は、読み出すべきデータがキャッシュメモリ2・2内に存在するか否かをチ

22 ックする(ステップ405、406)。

【0052】(6)ステップ406のチェックで、そのデータがキャッシュメモリ2・2内に存在した場合、そのデータをキャッシュメモリ2・2から読み出し、また、そ

23 のデータをキャッシュメモリ2・2から読み出し、また、そ

図13において、各論理記憶領域や物理記憶領域のアクセス占有時間を構成して、物理記憶装置やパリティグループ毎のアクセス占有時間と求めた値に基づいて新たに占有時間削除テーブルを作成してもよい。

【0064】前述により、ホスト10や別端末17等が物理記憶装置の移動性取扱要求をディスクアレイサ

ある論理記憶領域のアドレスとその論理記憶領域内の物理記憶領域のアドレスとの対応を管理している。これに使用する論理／物理対が情報 1-1 は、図 7 に示すように、特定の論理記憶領域に対して付与される論理記憶領域番号 510、論理アドレス 510、その論理記憶領域番号 500、論理アドレス 510 と、物理記憶領域番号 521 と物理記憶領域番号 521 と物

理外に記述要領のノドレス5 2 2 による物理アドレス  
5 2 0 、その物理配接装置1 5 の性能を示すレイドレバ  
ル5 3 0 、その物理配接装置1 5 が属しているバリティ  
グループ番号5 0 0 のそれを対応付けられて構成さ  
れる。ディスクアレイ制御装置1 4 は、このようにスト1 0  
ノードに対する情報1 4 1 を有することにより、ホスト1 0  
や制御端末1 1 7 、その他のチップユール(例えば、その他の  
物理配接装置1 4 1 を有する)との間でデータ交換を行  
なれば、データ書き込み終了通知をディスクアレイ制御装置1 4  
1 3 に通知すると共に、前述のタイトデータを物理記憶  
領域1 2 3 に格納する(ステップ4 1 2 、4 1 3)。

[0055] (9) ステップ4 0 8 の処理後、または、  
[0055] (9) ステップ4 0 8 の処理後、または、

節2.2内の構成・分類情報管理部2.2.6に格納される物理記憶領域のアドレスと物理記憶領域のアドレスとの対応付けを管理するテーブルの構成例を示す図である。

【図006】物理記憶域1～5内のデータ処理部12  
2.4は、ティスクアレイ前倒置表13より、ストレージサブシステム12の操作時や、物理記憶域の増減やヘッド交換時の移動用データベース（アーカイブ）データベースとの接続

40 のディスクアレイ制御装置等からリードやライト等の物理記憶領域のアドレスを指定して自ディスクアレイ制御装置1-3にアクセスされた場合、物理領域のアドレスを物理的な記憶領域のアドレスに変換してデータのリード/ライト処理を物理記憶装置1-5に對して正確に行なうことができる。

る論理記憶領域のアドレスと物理記憶領域のアドレスとの関係が変化した際に、その論理記憶領域と物理記憶領域との対応付け情報を受信し、それを構成・分類情報管理部226に図9に示すような状態テーブルに格納する。この状態テーブルは、特定の論理記憶領域に対して付与される論理記憶領域番号000、論理アドレス71

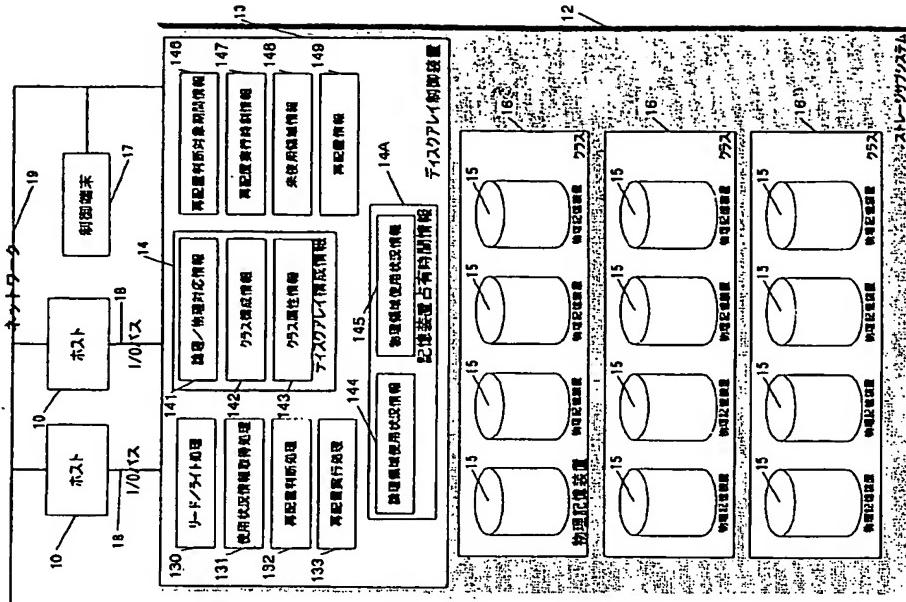
に、キャッシュモード 2 あるいは物理記憶領域 2.3 にアクセスした占有時間情報を錄音情報取扱部 2.2 内の占有時間処理部 2.7 に格納する(ステップ 4.1、4.2)。

い、この物理記憶アドレスを用いて、物理記憶アドレス722に登録された物理記憶装置721を主導する出力回路723が動作する。これにより、物理記憶アドレス722に登録された物理記憶装置721が動作する。これにより、物理記憶アドレス722に登録された物理記憶装置721が動作する。

- り、より最適なストレージシステムの性能のチューニングを行うことができる。
- 【図面の説明】
- 【図1】本発明によるストレージシステムを構成した計算機システムの構成を示すブロック図である。
- 【図2】物理記憶装置の構成を示すブロック図である。
- 【図3】ストレージシステムが起動されたときのディスクアレイ制御装置の処理動作を説明するフローチャートである。
- 【図4】物理記憶領域のアドレスと物理記憶領域のアドレスが変化した場合のディスクアレイ制御装置側の処理動作を説明するフローチャートである。
- 【図5】ディスクアレイ制御装置が物理記憶装置の移動情報取扱部の情報を読み出す際のディスクアレイ制御装置側の処理動作を説明するフローチャートである。
- 【図6】物理記憶装置内の物理記憶制御装置側の処理動作を説明するフローチャートである。
- 【図7】ディスクアレイ制御装置側に保持されている物理記憶領域とのアドレスの対応を管理する船積管理／物理対応情報の構成例を説明する図である。
- 【図8】ディスクアレイ制御装置側に格納される船積管理情報と物理記憶領域使用状況情報等の記憶装置占有時間情報の例を示す図である。
- 【図9】物理記憶装置内の構成・分担情報管理部に格納される船積管理情報と物理記憶領域とのアドレスの対応付けを管理するテーブルの構成例を示す図である。
- 【図10】物理記憶装置内の移動情報取扱部の占有時間統計部に保存・格納される記憶領域へのアクセスによる占有時間情報のテーブルの構成例を示す図である。
- 【符号の説明】
- 10 ポスト

【図1】

図1

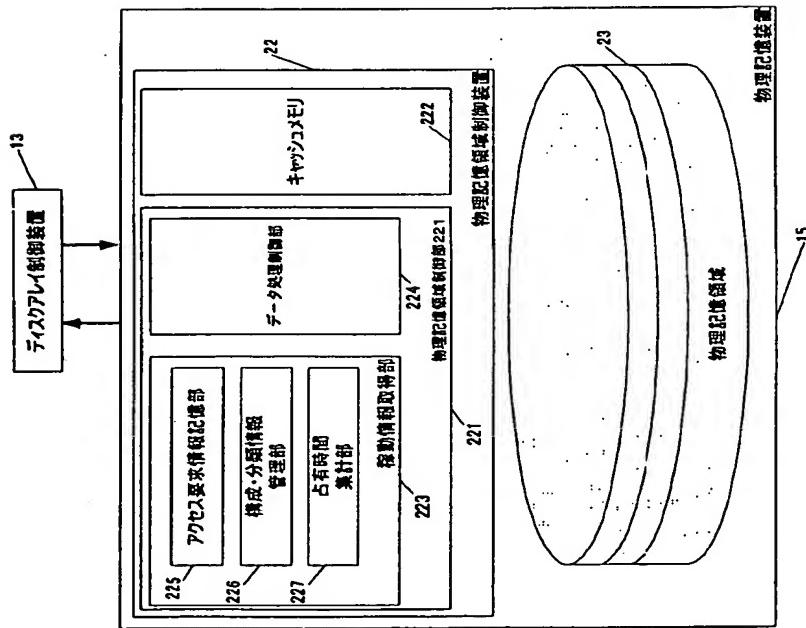
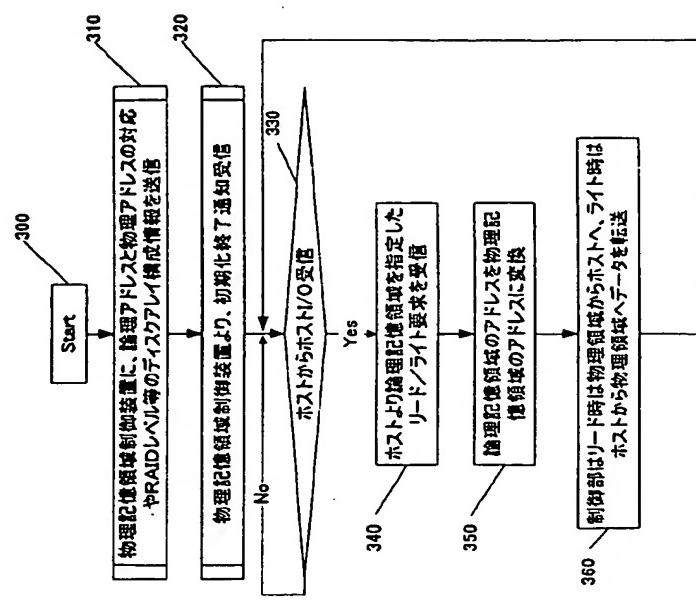


操作	V0	Sequential Read	Sequential Write	Random Read	Random Write	Random write Data Party	Random write Party	Cache	合計
1	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	252
2	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	163
3	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	163
4	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	占有時間	163
...	...	...	...	...	...	...	...	...	...

【図8】

操作	\$02	\$10	\$20	\$30	\$40	\$50	\$60	\$70	\$80
1	占有時間								
2	占有時間								
3	占有時間								
4	占有時間								
...	...	...	...	...	...	...	...	...	...

【図8】

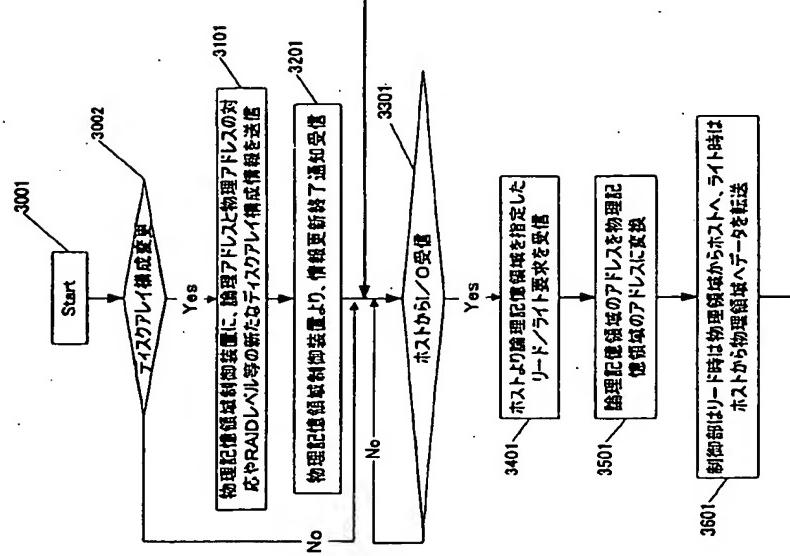
【図2】  
図2【図3】  
図3

論理記憶領域番号	論理アドレス	物理記憶装置番号	物理記憶領域アドレス
3	3000~3999	1	0~999
4	4000~4999	1	1000~1999
5	5000~5999	1	2000~2999

【図9】  
図9

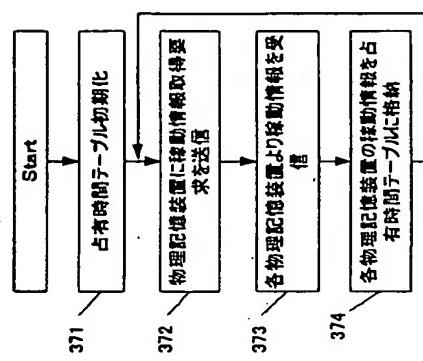
【図4】

図4



【図5】

図5



【図6】

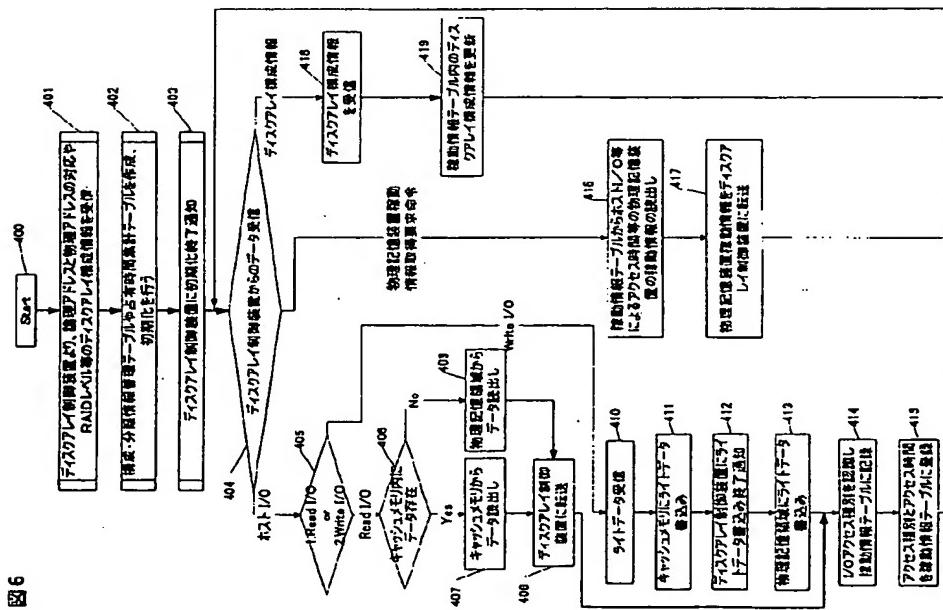
図6

物理記憶装置番号	物理アドレス	物理記憶装置番号 物理アドレス	レイドレベル	バリティグループ番号
0	0~99	0 ~99	0~199	1 100
1	1000~1999	0	1000~1999	1 100
2	2000~2999	0	2000~2999	1 100
3	3000~3999	1	0~999	5 120
4	4000~4999	1	1000~1999	5 120
5	5000~5999	1	2000~2999	5 120
...	...	...	...	...

【図7】

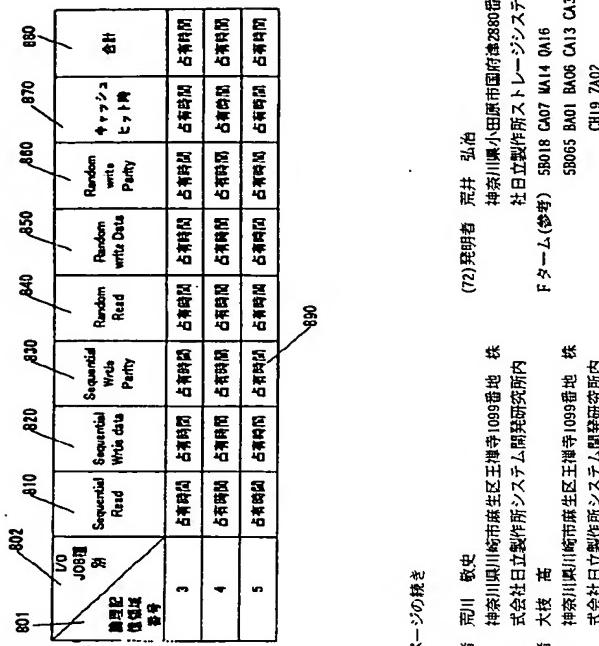
図7

【図6】



【図6】

【図10】



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